

# ECON2125/8013

## Lecture 3

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# Announcements, Reminders

## 1. Tutorials

- Start this week
- Questions are on-line
- Overcrowding

## 2. Revision / weak prerequisites

- Can audit EMET1001

## 3. Office hours

- Apologies

# Constrained Optimization

A major focus of econ: the optimal allocation of scarce resources

- Optimal means optimization, scarce means constrained

Standard constrained problems:

- Maximize utility given budget
- Maximize portfolio return given risk constraints
- Minimize cost given output requirement

**Example.** Maximization of utility subject to budget constraint

$$\max u(x_1, x_2) \text{ s.t. } p_1x_1 + p_2x_2 \leq m$$

Here

- $p_i$  is the price of good  $i$ , assumed  $> 0$
- $m$  is the budget, assumed  $> 0$
- $x_i \geq 0$  for  $i = 1, 2$

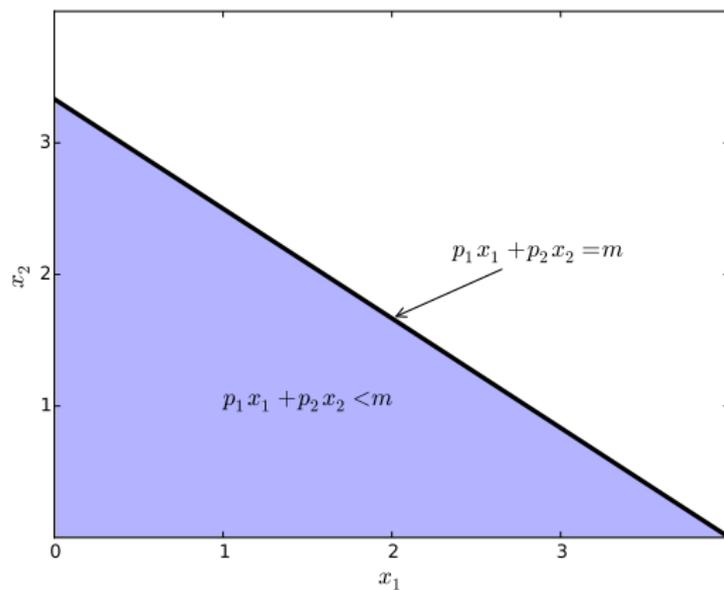


Figure : Budget set when,  $p_1 = 1$ ,  $p_2 = 1.2$ ,  $m = 4$

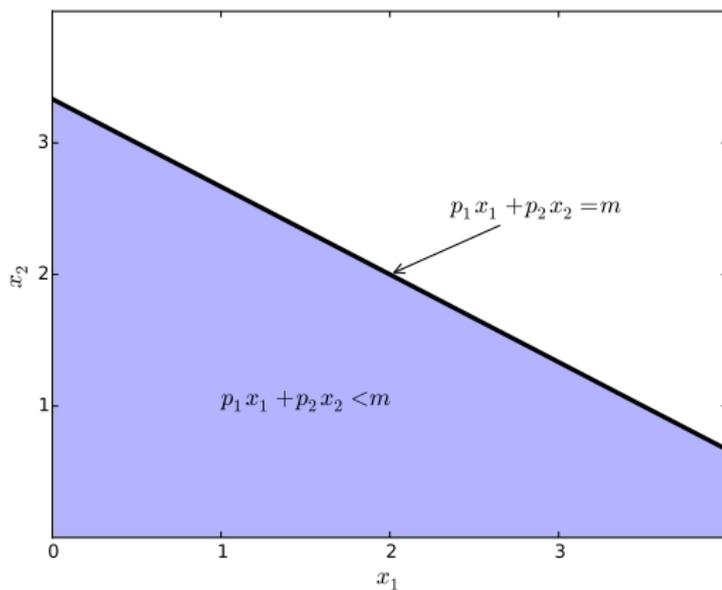


Figure : Budget set when,  $p_1 = 0.8$ ,  $p_2 = 1.2$ ,  $m = 4$

**Example.** Suppose we want to solve

$$\max u(x_1, x_2) \quad \text{s.t.} \quad p_1x_1 + p_2x_2 \leq m$$

Let's assume that

$$u(x_1, x_2) = \alpha \log(x_1) + \beta \log(x_2)$$

where

- $0 < \alpha, \beta$

Let's recall the utility function's shape

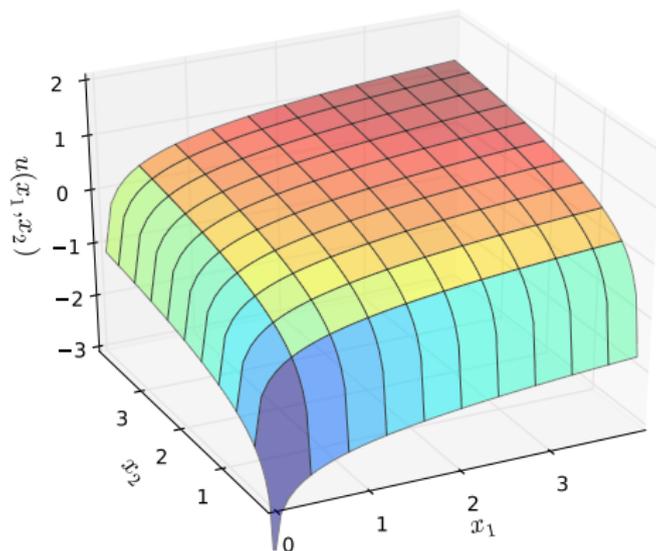


Figure : Log utility with  $\alpha = 0.4$ ,  $\beta = 0.5$

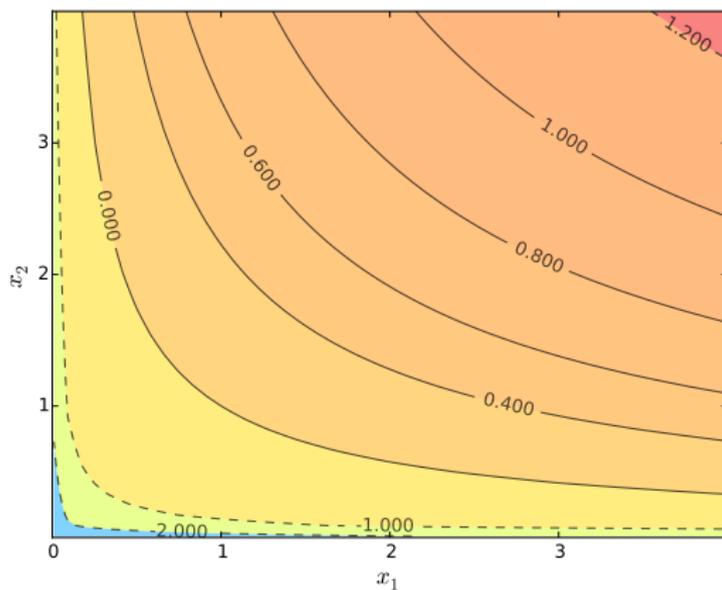


Figure : Log utility with  $\alpha = 0.4$ ,  $\beta = 0.5$

We seek a bundle  $(x_1^*, x_2^*)$  that maximizes  $u$  over the budget set

That is,

$$\alpha \log(x_1^*) + \beta \log(x_2^*) \geq \alpha \log(x_1) + \beta \log(x_2)$$

for all  $(x_1, x_2)$  satisfying  $x_i \geq 0$  for each  $i$  and

$$p_1 x_1 + p_2 x_2 \leq m$$

Visually, here is the budget set and objective function:

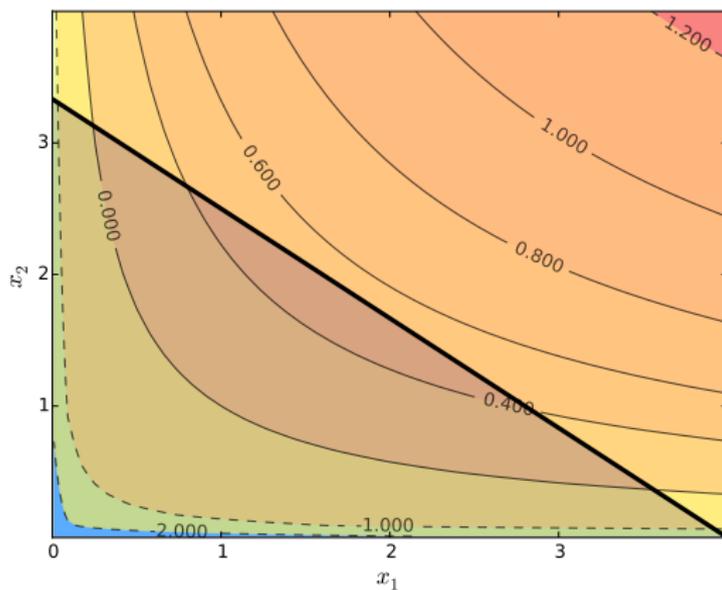


Figure : Utility max for  $p_1 = 1$ ,  $p_2 = 1.2$ ,  $m = 4$ ,  $\alpha = 0.4$ ,  $\beta = 0.5$

First observation:  $u(0, x_2) = u(x_1, 0) = u(0, 0) = -\infty$

- Hence we need consider only strictly positive bundles

Second observation:  $u(x_1, x_2)$  is strictly increasing in both  $x_i$

- Never choose a point  $(x_1, x_2)$  with  $p_1x_1 + p_2x_2 < m$
- Otherwise can increase  $u(x_1, x_2)$  by small change

Hence we can replace  $\leq$  with  $=$  in the constraint

$$p_1x_1 + p_2x_2 \leq m \quad \text{becomes} \quad p_1x_1 + p_2x_2 = m$$

Implication: Just search along the budget line

# Substitution Method

Substitute constraint into objective function

This changes

$$\max_{x_1, x_2} \{ \alpha \log(x_1) + \beta \log(x_2) \} \text{ s.t. } p_1 x_1 + p_2 x_2 = m$$

into

$$\max_{x_1} \{ \alpha \log(x_1) + \beta \log((m - p_1 x_1) / p_2) \}$$

Since all candidates satisfy  $x_1 > 0$  and  $x_2 > 0$ , the domain is

$$0 < x_1 < \frac{m}{p_1}$$

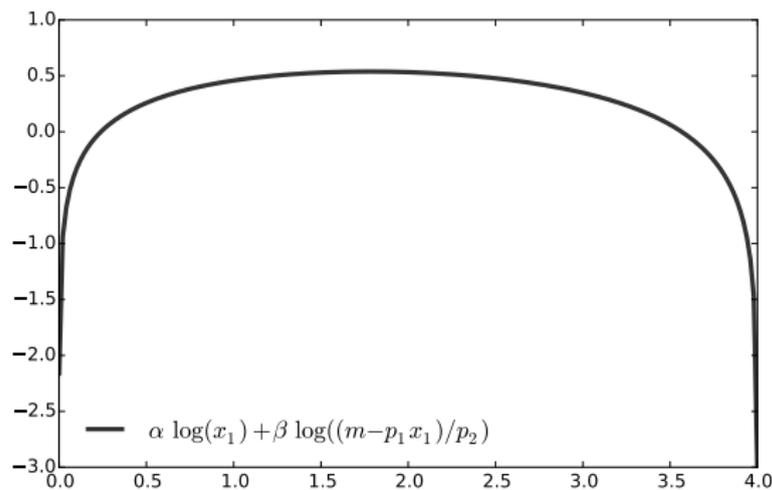


Figure : Utility max for  $p_1 = 1$ ,  $p_2 = 1.2$ ,  $m = 4$ ,  $\alpha = 0.4$ ,  $\beta = 0.5$

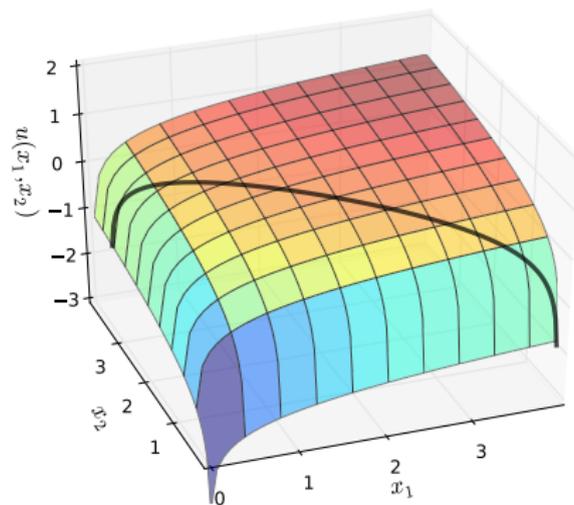


Figure : Utility max for  $p_1 = 1$ ,  $p_2 = 1.2$ ,  $m = 4$ ,  $\alpha = 0.4$ ,  $\beta = 0.5$

First order condition for

$$\max_{x_1} \{ \alpha \log(x_1) + \beta \log((m - p_1 x_1) / p_2) \}$$

gives

$$x_1^* = \frac{\alpha}{\alpha + \beta} \cdot \frac{m}{p_1}$$

**Ex.** Verify

**Ex.** Check second order condition (strict concavity)

Substituting into  $p_1 x_1^* + p_2 x_2^* = m$  gives

$$x_2^* = \frac{\beta}{\beta + \alpha} \cdot \frac{m}{p_2}$$

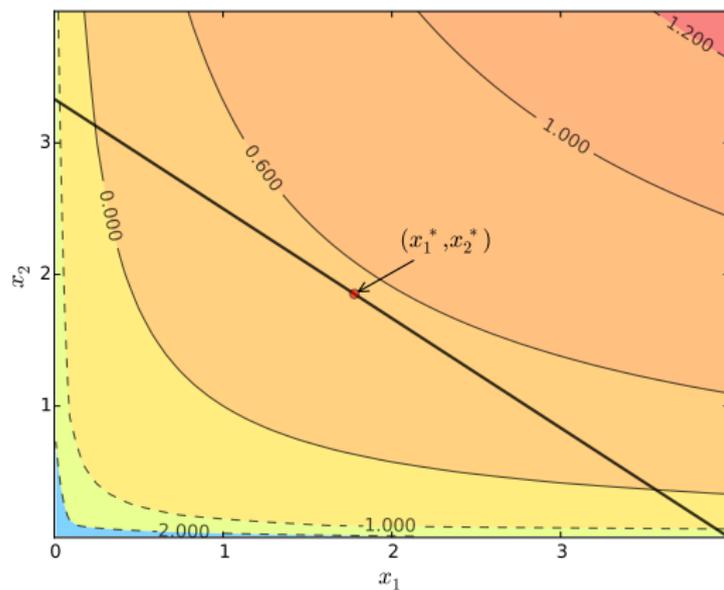


Figure : Maximizer for  $p_1 = 1$ ,  $p_2 = 1.2$ ,  $m = 4$ ,  $\alpha = 0.4$ ,  $\beta = 0.5$

# Substitution Method Cookbook

How to solve

$$\max_{x_1, x_2} f(x_1, x_2)$$

$$\text{s.t. } g(x_1, x_2) = 0$$

Steps:

1. Write constraint as  $x_2 = h(x_1)$  for some function  $h$
2. Solve univariate problem  $\max_{x_1} f(x_1, h(x_1))$  to get  $x_1^*$
3. Plug  $x_1^*$  into  $x_2 = h(x_1)$  to get  $x_2^*$

### Example. (Minimization)

Consider the simple problem

$$\min_{x_1, x_2} \{x_1^2 + x_2^2\}$$

$$\text{s.t. } x_1 + x_2 - 10 = 0$$

1. Write constraint as  $x_2 = 10 - x_1$
2. Solve  $\min_{x_1} \{x_1^2 + (10 - x_1)^2\}$  to get  $x_1^* = 5$
3. Plug  $x_1^* = 5$  into  $x_1 + x_2 = 10$  to get  $x_2^* = 5$

## Limitations

Substitution doesn't always work

**Example.** Suppose that  $g(x_1, x_2) = x_1^2 + x_2^2 - 1$

Step 1 from the cookbook says

use  $g(x_1, x_2) = 0$  to write  $x_2$  as a function of  $x_1$

But  $x_2$  has two possible values for each  $x_1 \in (-1, 1)$

$$x_2 = \pm \sqrt{1 - x_1^2}$$

In other words,  $x_2$  is not a well defined function of  $x_1$

As we meet more complicated constraints such problems intensify

- Constraint defines complex curve in  $(x_1, x_2)$  space
- Inequality constraints, etc.

We need more general, systematic approaches too

Leads to next discussion

# Tangency

Consider again the problem

$$\max_{x_1, x_2} \{ \alpha \log(x_1) + \beta \log(x_2) \}$$

$$\text{s.t. } p_1 x_1 + p_2 x_2 = m$$

Turns out that the maximizer has the following property:

- budget line is tangent to an indifference curve at maximizer

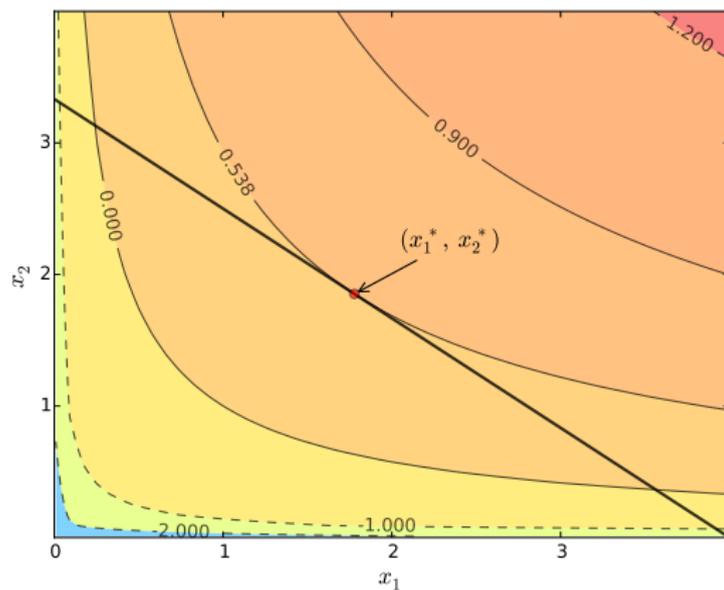


Figure : Maximizer for  $p_1 = 1$ ,  $p_2 = 1.2$ ,  $m = 4$ ,  $\alpha = 0.4$ ,  $\beta = 0.5$

In fact this is an instance of a general pattern

Notation: Let's call  $(x_1, x_2)$  interior to the budget line if  $x_i > 0$  for  $i = 1, 2$

- Not a “corner” solution

In general, any interior maximizer  $(x_1^*, x_2^*)$  of differentiable utility function  $u$  has the property

budget line is tangent to a contour line at  $(x_1^*, x_2^*)$

Otherwise we can do better:

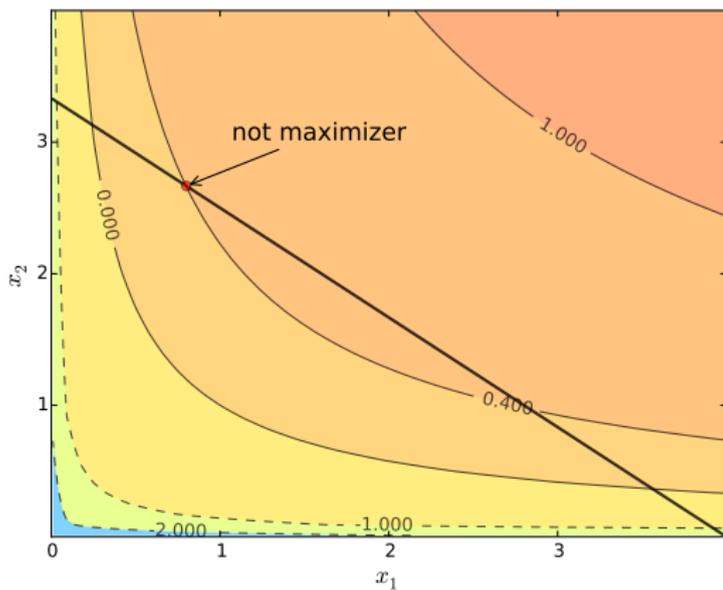


Figure : When tangency fails we can do better

Necessity of tangency often rules out a lot of points

Can we exploit this fact to

- Build intuition
- Develop more general methods?

The answer is yes

## Using Tangency: Relative Slope Conditions

- Relies on tangency idea discussed above
- Generalizes nicely

Consider the smooth, equality constrained optimization problem

$$\max_{x_1, x_2} f(x_1, x_2)$$

$$\text{s.t. } g(x_1, x_2) = 0$$

How to develop necessary conditions for optima via tangency?

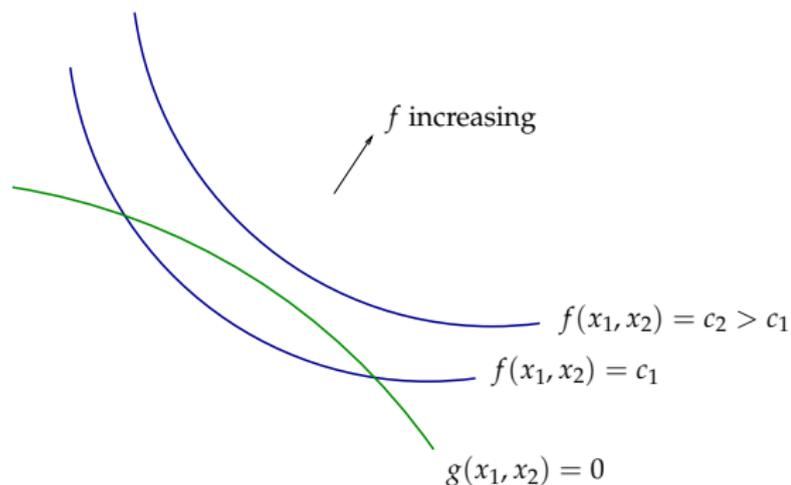


Figure : Contours for  $f$  and  $g$

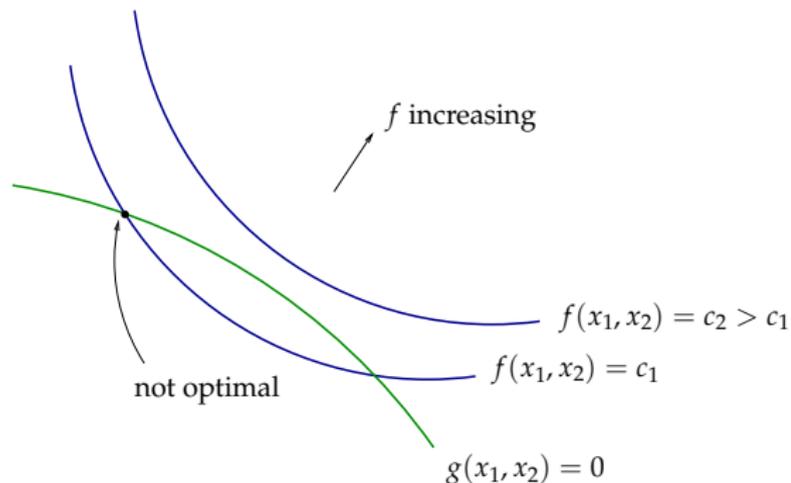


Figure : Contours for  $f$  and  $g$

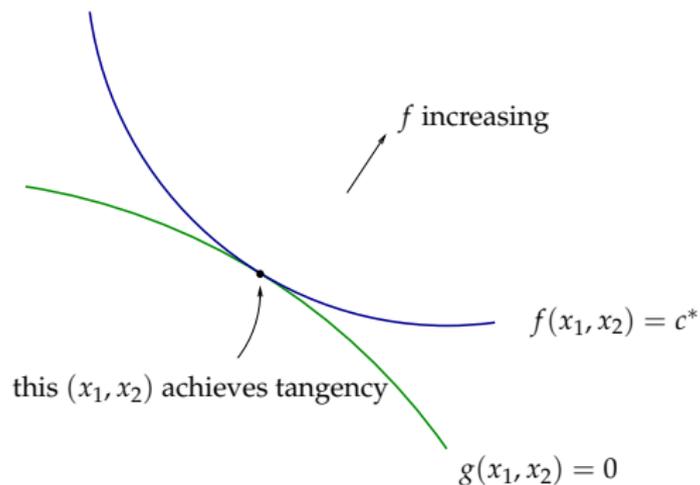


Figure : Tangency necessary for optimality

How do we locate such an  $(x_1, x_2)$  pair?

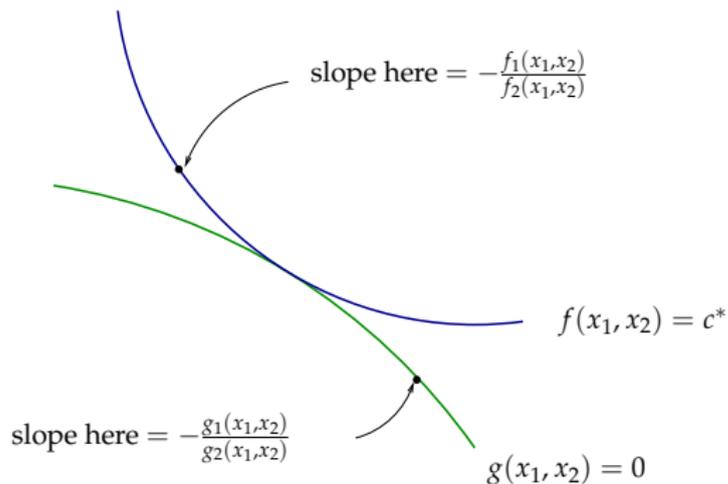


Figure : Slope of contour lines

Sketch of proof for case of  $f$

Let's fix  $c$  and vary  $x_2$  with  $x_1$  to maintain  $f(x_1, x_2) = c$

This implicitly defines  $x_2$  as a function of  $x_1$

The slope of this function is what we're after

Differentiating  $f(x_1, x_2(x_1)) = c$  with respect to  $x_1$  gives

$$f_1(x_1, x_2) + f_2(x_1, x_2)x_2'(x_1) = 0$$

Solving gives slope  $= x_2'(x_1) = -f_1(x_1, x_2) / f_2(x_1, x_2)$

Proper proof: See formula for implicit differentiation

Now let's choose  $(x_1, x_2)$  to equalize the slopes

That is, choose  $(x_1, x_2)$  to solve

$$-\frac{f_1(x_1, x_2)}{f_2(x_1, x_2)} = -\frac{g_1(x_1, x_2)}{g_2(x_1, x_2)}$$

Equivalent:

$$\frac{f_1(x_1, x_2)}{f_2(x_1, x_2)} = \frac{g_1(x_1, x_2)}{g_2(x_1, x_2)}$$

Also need to respect  $g(x_1, x_2) = 0$

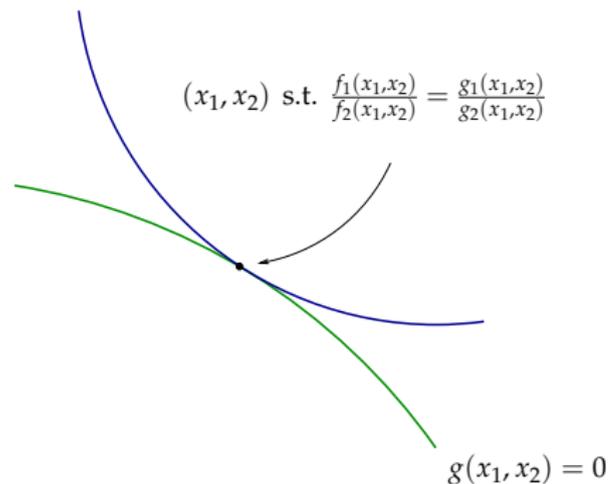


Figure : Condition for tangency

# Tangency Condition Cookbook

In summary, when  $f$  and  $g$  are both differentiable functions, to find candidates for optima in

$$\begin{aligned} \max_{x_1, x_2} f(x_1, x_2) \\ \text{s.t. } g(x_1, x_2) = 0 \end{aligned}$$

1. (Impose slope tangency) Set

$$\frac{f_1(x_1, x_2)}{f_2(x_1, x_2)} = \frac{g_1(x_1, x_2)}{g_2(x_1, x_2)}$$

2. (Impose constraint) Set  $g(x_1, x_2) = 0$

3. Solve simultaneously for  $(x_1, x_2)$  pairs satisfying these conditions

Example. Consider again

$$\begin{aligned} \max_{x_1, x_2} \{ & \alpha \log(x_1) + \beta \log(x_2) \} \\ \text{s.t. } & p_1 x_1 + p_2 x_2 - m = 0 \end{aligned}$$

Then

$$\frac{f_1(x_1, x_2)}{f_2(x_1, x_2)} = \frac{g_1(x_1, x_2)}{g_2(x_1, x_2)} \iff \frac{\alpha x_2}{\beta x_1} = \frac{p_1}{p_2}$$

Solving simultaneously with  $p_1 x_1 + p_2 x_2 = m$  gives

$$x_1^* = \frac{\alpha}{\alpha + \beta} \cdot \frac{m}{p_1} \quad \text{and} \quad x_2^* = \frac{\beta}{\beta + \alpha} \cdot \frac{m}{p_2}$$

Same as before...

# Slope Conditions for Minimization

Good news: The conditions are exactly the same

In particular:

- Lack of tangency means not optimizer
- Constraint must be satisfied

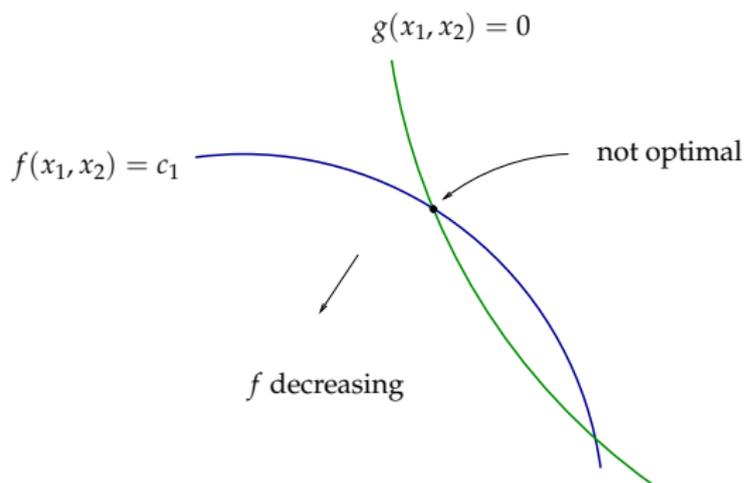


Figure : Lack of tangency

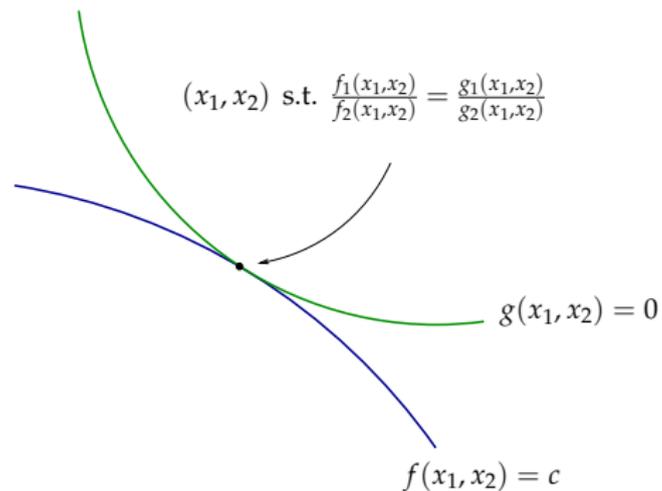


Figure : Condition for tangency

**Example.** Minimize cost for a given level of production  $q$

$$\begin{aligned} \min_{k,\ell} & \{rk + w\ell\} \\ \text{s.t.} & Ak^\alpha \ell^\beta \geq q \end{aligned}$$

All parameters assumed to be strictly positive

Since inputs are costly, any minimizer will produce exactly  $q$

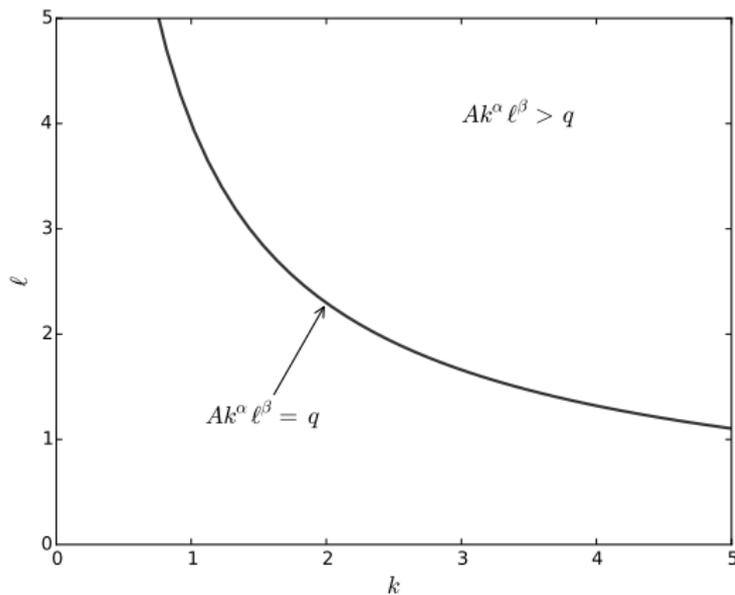


Figure : Constraint set when  $A = 2.0$ ,  $\alpha = 0.4$ ,  $\beta = 0.5$ ,  $q = 4.0$

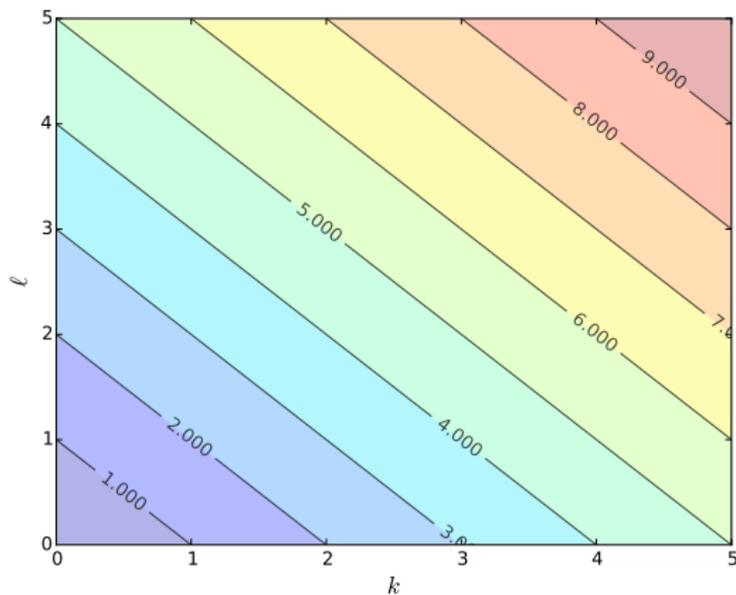


Figure : Objective function  $rk + wl$  when  $w = r = 1.0$

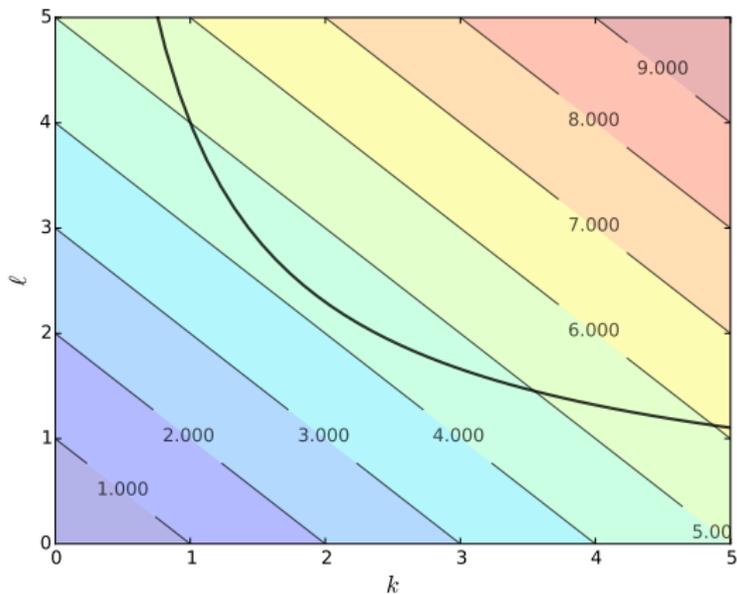


Figure : Together, same parameters

To apply the method we set

$$g(k, \ell) := Ak^\alpha \ell^\beta - q = 0$$

and

$$f(k, \ell) = rk + w\ell$$

Now let's apply the conditions to obtain candidates for minima

Slope tangency condition:

$$\frac{f_1(k, \ell)}{f_2(k, \ell)} = \frac{g_1(k, \ell)}{g_2(k, \ell)} \iff \frac{r}{w} = \frac{A\alpha k^{\alpha-1} \ell^\beta}{A k^\alpha \beta \ell^{\beta-1}} = \frac{\alpha \ell}{\beta k}$$

Combine this with the constraint

$$g(k, \ell) = Ak^\alpha \ell^\beta - q = 0$$

to get

$$k^* = \left[ \frac{q}{A} \left( \frac{w\alpha}{r\beta} \right)^\beta \right]^{1/(\alpha+\beta)} \quad \text{and} \quad \ell^* = \left[ \frac{q}{A} \left( \frac{r\beta}{w\alpha} \right)^\alpha \right]^{1/(\alpha+\beta)}$$

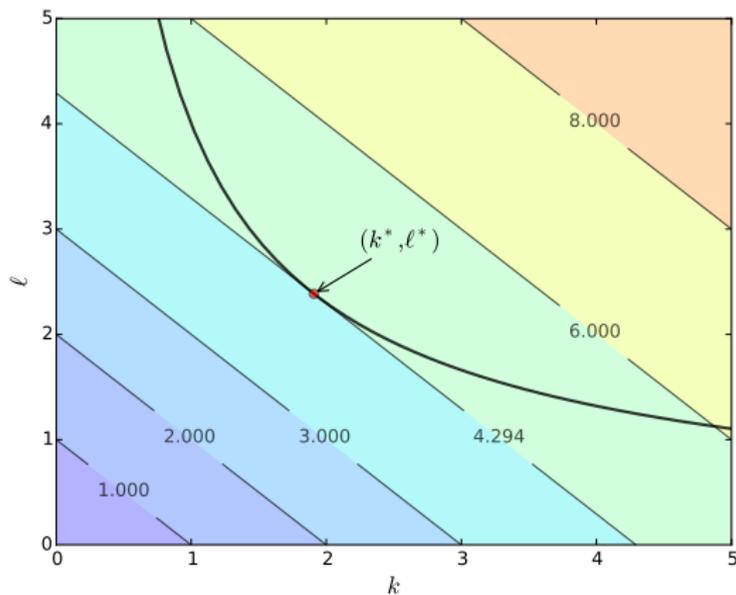


Figure : Minimizer  $(k^*, l^*)$

### Example. Intertemporal problem

$$\begin{aligned} \max_{c_1, c_2} U(c_1, c_2) &:= u(c_1) + \beta u(c_2) \\ \text{s.t. } c_2 &\leq (1+r)(w - c_1) \end{aligned}$$

where

- $r$  = interest rate,  $w$  = wealth,  $\beta$  = discount factor
- all parameters  $> 0$  and  $u$  strictly increasing

Write constraint as  $g(c_1, c_2) := (1+r)(w - c_1) - c_2 = 0$

Any interior solution must satisfy tangency condition

$$\frac{U_1(c_1, c_2)}{U_2(c_1, c_2)} = \frac{g_1(c_1, c_2)}{g_2(c_1, c_2)} \iff \frac{u'(c_1)}{u'(c_2)} = \beta(1+r)$$

## Method of Lagrange

The “standard machine” for optimization with equality constraints

$$\max_{x_1, x_2} f(x_1, x_2) \quad \text{s.t.} \quad g(x_1, x_2) = 0$$

Set

$$\mathcal{L}(x_1, x_2, \lambda) := f(x_1, x_2) + \lambda g(x_1, x_2)$$

and solve

$$\frac{\partial}{\partial x_1} \mathcal{L} = 0, \quad \frac{\partial}{\partial x_2} \mathcal{L} = 0, \quad \frac{\partial}{\partial \lambda} \mathcal{L} = 0$$

simultaneously

Since  $\mathcal{L}(x_1, x_2, \lambda) := f(x_1, x_2) + \lambda g(x_1, x_2)$  we have

$$\frac{\partial}{\partial x_i} \mathcal{L}(x_1, x_2, \lambda) = 0 \iff f_i(x_1, x_2) = -\lambda g_i(x_1, x_2), \quad i = 1, 2$$

Hence  $\frac{\partial}{\partial x_1} \mathcal{L} = \frac{\partial}{\partial x_2} \mathcal{L} = 0$  gives

$$\frac{f_1(x_1, x_2)}{f_2(x_1, x_2)} = \frac{g_1(x_1, x_2)}{g_2(x_1, x_2)}$$

Finally

$$\frac{\partial}{\partial \lambda} \mathcal{L}(x_1, x_2, \lambda) = 0 \iff g(x_1, x_2) = 0$$

Hence the method leads us to the same conditions

# Extensions

Let's look at some problems and extensions

Remark 1: The direct tangent slope condition can fail if we're dividing by zero in

$$\frac{f_1(x_1, x_2)}{f_2(x_1, x_2)} = \frac{g_1(x_1, x_2)}{g_2(x_1, x_2)}$$

In this case try the more general Lagrange conditions

$$f_1(x_1, x_2) + \lambda g_1(x_1, x_2) = 0$$

$$f_2(x_1, x_2) + \lambda g_2(x_1, x_2) = 0$$

Remark 2: Consider the two optimization problems

$$\max_{x_1, x_2} \{ \alpha \log(x_1) + \beta \log(x_2) \}$$

$$\text{s.t. } p_1 x_1 + p_2 x_2 = m$$

and

$$\max_{x_1, x_2} x_1^\alpha x_2^\beta$$

$$\text{s.t. } p_1 x_1 + p_2 x_2 = m$$

The tangency conditions are identical

$$\frac{\alpha x_2}{\beta x_1} = \frac{p_1}{p_2} \quad \text{and} \quad = \frac{\alpha x_1^{\alpha-1} x_2^\beta}{\beta x_1^\alpha x_2^{\beta-1}} = \frac{\alpha x_2}{\beta x_1} = \frac{p_1}{p_2}$$

More generally, maximizers are unchanged by increasing transformations

- Can be useful to simplify your problem

On the other hand maximum values are changed, of course

More on this later...

## Corner Solutions

So far all our solutions have been interior ( $x_i > 0$  for  $i = 1, 2$ )

Such solutions can be tracked down by the tangency conditions

However sometimes solutions are naturally on the boundaries

**Example.** Maximize  $x_1 + \log(x_2)$  subject to

$$p_1x_1 + p_2x_2 = m \quad \text{and} \quad x_1, x_2 \geq 0$$

Let's try the tangency approach with  $p_1 = p_2 = 1$  and  $m = 0.4$

Tangency condition is

$$\frac{1}{1/x_2} = \frac{p_1}{p_2} \iff x_2 = \frac{p_1}{p_2} = 1$$

Applying the budget constraint gives

$$x_1 + x_2 = 0.4 \quad \text{and hence} \quad x_1 = -0.6$$

Meaning: There is no tangent point in

$$D := \{(x_1, x_2) : x_1 \geq 0, x_2 \geq 0 \text{ and } p_1x_1 + p_2x_2 = m\}$$

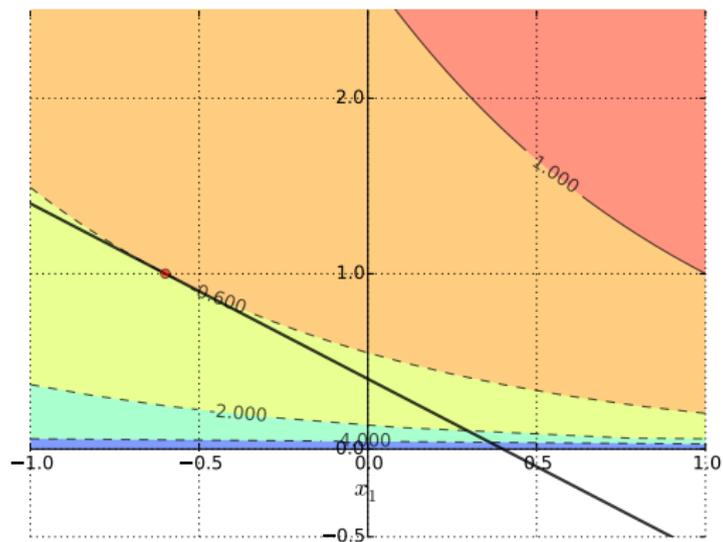


Figure : Tangent point is infeasible

Interpretation: No interior solution

Put differently

- At every interior point on the budget line you can do better
- Hence solution must be on the boundary

Since  $x_2 = 0$  implies  $x_1 + \log(x_2) = -\infty$ , solution is

- $x_1^* = 0$
- $x_2^* = m/p_2 = 0.4$

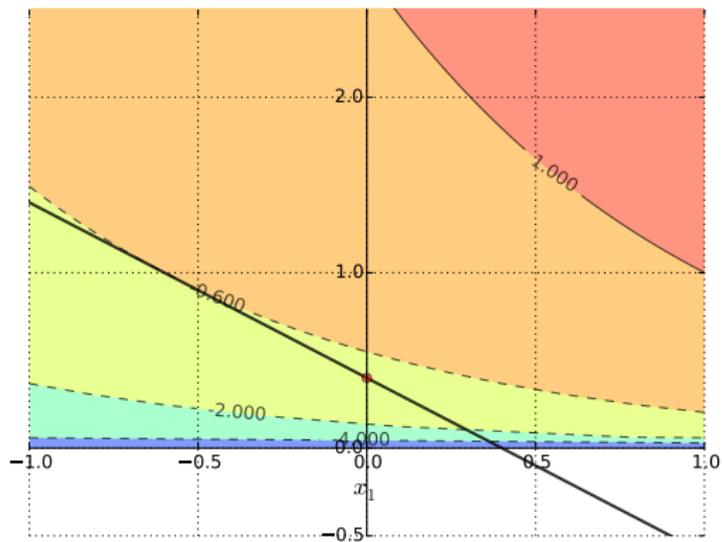


Figure : Corner solution